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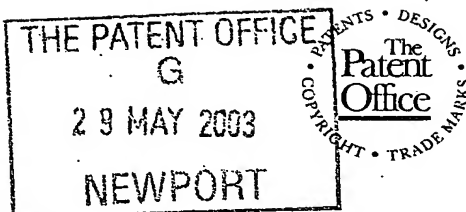
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3. Full name, address and postcode of the or of each applicant (underline all surnames)

Hewlett-Packard Development Company, L.P.
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USA

Patents ADP number (if you know it)

8557886001

If the applicant is a corporate body, give the country/state of its incorporation

4. Title of the invention Secure Data Provision Method and Apparatus and Data Recovery Method and System

5. Name of your agent (if you have one)

"Address for service" in the United Kingdom to which all correspondence should be sent (including the postcode)

Robert F. Squibbs
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Bristol BS34 8QZ

Patents ADP number (if you know it)

7928187001

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Country

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0308891.1

17/04/2003

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Robert F. Squibbs
Robert F. Squibbs

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12. Name and daytime telephone number of person to contact in the United Kingdom

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Secure Data Provision Method and Apparatus and Data Recovery Method and System.

5 Field of the Invention

The present invention relates to a method and apparatus for the provision of target data in encrypted form to an accredited professional and to a method and system for recovering the target data in clear; in particular, but not exclusively, the present invention relates to such methods, system and apparatus involving Identifier-Based Encryption.

10

As used herein, reference to a "professional" is a reference to an individual that has certain recognised skills that the individual uses in carrying out their job. Such skills may range from the skills of a brain surgeon to those of a plumber or the like, without limitation.

15 Background of the Invention

Professionals working in the same field frequently belong to a professional body one role of which may be to maintain a list of accredited professionals working in the field concerned (though not necessarily members of the body); such a role may, indeed, have regulatory force. Entry on the list of accredited professionals often requires an individual to have obtained certain qualifications but will generally also require that the individual has not committed any major act detrimental to their clients. Thus the accredited status of a professional is not something which once obtained will necessarily continue.

One field where the professional status of an individual is of particular importance is the medical field. This field places high demands not only on the skill of the individuals concerned but also on maintaining the confidentiality of patient records. It is expected that electronic medical records of patients will replace paper records in the near future. The update of these records is likely to be the responsibility of the patient's local doctor (that is, their "general practitioner" or "GP"). The GP, for the purpose of secure preservation of patient data, is likely to use a secure data storage service to store the electronic patient records. In an emergency situation, in which a patient requires medical care, an attending doctor or paramedic (generally, a medical professional) needs to know, as a matter of

urgency, the medical history of the patient to prevent giving inappropriate treatments. There is therefore a need for the attending medical professional to obtain the patient's medical records from the data storage service provider; however, this needs to be done in a manner that safeguards the privacy of the records.

5

Most solutions that have been proposed for dealing with the above situation involve the use of a public key infrastructure (PKI) which would need to be created for the medical professionals. In such a PKI, a professional body for medical professionals would act as a certificate authority providing an accredited medical professional with a certificate
10 confirming their accreditation and public key. In an emergency situation, the medical professional would send a patient identifier together with the professional's own certificate to the patient data storage service. This service would verify the validity of the certificate, encrypt the patient's records with the medical professional's public key, and return the encrypted data to the medical professional.

15

One disadvantage of the foregoing arrangement is that it does not distinguish between a request from a medical professional carrying out their work in a hospital emergency room and a medical professional who just wants to pry into the details of a patient. Another disadvantage is the need for the data storage service to keep, or have immediate access to,
20 an up-to-date certificate revocation list.

It is an object of the present invention to provide an improved way for professionals to access confidential data in a controlled manner that obviates at least some of the problems associated with prior systems. It is to be understood that the present invention is not limited
25 to the provision of sensitive data to medical professionals but is applicable to all types of professionals.

As will explained hereinafter, the preferred embodiments of the invention utilise Identifier-Based Encryption (IBE) which is an emerging cryptographic schema. For convenience, this
30 known schema will next be described with reference to Figure 1 of the accompanying drawings. In an IBE schema, a data provider 10 encrypts payload data 13 using an encryption key string 14 and public data 15 provided by a trusted authority 12; the data provider 10 then provides the encrypted payload data <13> to a recipient 11 who decrypts

it using a decryption key 16 provided by the trusted authority together with the latter's public data. The trusted authority's public data is derived by the authority using private data 17 and a one-way function 18. Important features of the IBE schema are that any kind of string (including a name, a role, etc.) can be used as an encryption key string 14, and
 5 that the generation of the decryption key 16 is effected by the trusted authority (process 19) using the encryption key string 14 and its private data 17, enabling the generation of the decryption key 16 to be postponed until needed for decryption. Because the encryption key string frequently contains data identifying the intended recipient (for example, by a required characteristic), the encryption key string is also known as the identifier string.

10

The public data 15 and private data 17 effectively form a base key pair which in any particular application is typically only changed infrequently (if at all) as compared to the frequency of change of the encryption / decryption key pair 14,16.

15 A number of IBE algorithms are known, one of which is the "Quadratic Residuosity" (QR) method described in the paper: C. Cocks, "An identity based encryption scheme based on quadratic residues", Proceedings of the 8th IMA International Conference on Cryptography and Coding, LNCS 2260, pp 360-363, Springer-Verlag, 2001. A brief description of this form of IBE is given below.

20

In the QR method, the trusted authority's public data 15 comprises a value N that is a product of two random prime numbers p and q , where the values of p and q are the private data 17 of the trusted authority 12. The values of p and q should ideally be in the range of 2^{511} and 2^{512} and should both satisfy the equation: $p, q \equiv 3 \pmod{4}$. However, p and q must
 25 not have the same value. Also provided is a hash function $\#$ which when applied to a string returns a value in the range 0 to $N-1$.

Each bit m of the user's payload data 13 is then encrypted as follows:

- The data provider 10 generates random numbers t_+ (where t_+ is an integer in the range $[0, 2^N]$) until a value of t_+ is found that satisfies the equation $jacobi(t_+, N) = m$,
 30 where m has a value of -1 or 1 depending on whether the corresponding bit of the

user's data is 0 or 1 respectively. (As is well known, the *jacobi* function is such that where $x^2 \equiv \# \pmod{N}$ the *jacobi* ($\#, N$) = -1 if x does not exist, and = 1 if x does exist). The data provider 10 then computes the value:

$$s_+ \equiv (t_+ + \#(\text{encryption_keystring})/t_+) \pmod{N}$$

5 where s_+ corresponds to the encrypted value of the bit m concerned.

- Since $\#(\text{encryption_keystring})$ may be non-square, the data provider additionally generates additional random numbers t_- (integers in the range $[0, 2^N)$) until one is found that satisfies the equation $\text{jacobi}(t_-, N) = m$. The data provider 10 then computes the value:

$$s_- \equiv (t_- - \#(\text{encryption_keystring})/t_-) \pmod{N}$$

as the encrypted value of the bit m concerned.

The encrypted values s_+ and s_- for each bit m of the user's data are then made available to the intended recipient 11, for example via e-mail or by being placed in a electronic public area; the identity of the trusted authority 12 and the encryption key string 14 will generally also be made available in the same way.

The encryption key string 14 is passed to the trusted authority 12 by any suitable means; for example, the recipient 11 may pass it to the trusted authority or some other route is used - indeed, the trusted authority may have initially provided the encryption key string. The trusted authority 12 determines the associated private key B by solving the equation:

$$B^2 \equiv \#(\text{encryption_keystring}) \pmod{N} \quad (\text{"positive" solution})$$

If a value of B does not exist, then there is a value of B that is satisfied by the equation:

$$25 \quad B^2 \equiv -\#(\text{encryption_keystring}) \pmod{N} \quad (\text{"negative" solution})$$

As N is a product of two prime numbers p, q it would be extremely difficult for any one to calculate the decryption key B with only knowledge of the encryption key string and N . However, as the trusted authority 12 has knowledge of p and q (i.e. two prime numbers) it is relatively straightforward for the trusted authority 12 to calculate B .

Any change to the encryption key string 14 will result in a decryption key 16 that will not decrypt the payload data 13 correctly. Therefore, the intended recipient 11 cannot alter the encryption key string before supplying it to the trusted authority 12.

- 5 The trusted authority 12 sends the decryption key to the data recipient 11 along with an indication of whether this is the “positive” or “negative” solution for B.

If the “positive” solution for the decryption key has been provided, the recipient 11 can now recover each bit m of the payload data 13 using:

10
$$m = \text{jacobi}(s_+ + 2B, N)$$

If the “negative” solution for the decryption key B has been provided, the recipient 11 recovers each bit m using:

$$m = \text{jacobi}(s_- + 2B, N)$$

- 15 Whilst in the foregoing example, the encryption key string has been used directly in the QR IBE algorithm, it is also possible to use in the encryption process a derivative of the encryption key string, this derivative being formed, for example, by using a predetermined hash function. In this case, the entity generating the decryption key can still simply be supplied with the encryption key string provided it knows the predetermined function used
- 20 to form the derivative of the encryption key string (in fact, this is equivalent to using a variant of stated the QR IBE algorithm in which the predetermined function is applied to the encryption key string wherever the latter appears). Where the decryption-key generating entity does not need to access the contents of the original encryption key string, then it need only be provided with the derivative of the encryption key string used during the
- 25 encryption process. In the following description, where the term “encryption key” is used, this is intended to refer to the form of the encryption key string used in the stated version of IBE algorithm concerned whether this is the unprocessed encryption key string or a derivative formed by subjecting the encryption key string to predetermined processing.

- 30 Other IBE algorithms are known such as the use of Weil or Tate pairings – see, for example: D. Boneh, M. Franklin – “Identity-based Encryption from the Weil Pairing” in

Advances in Cryptology - CRYPTO 2001, LNCS 2139, pp. 213-229, Springer-Verlag, 2001. IBE algorithms based on the Weil or Tate pairings are usually described in terms of there being an IBE encryption key that is derived in a predetermined manner from an encryption key string (though it would be possible to re-state the algorithms such that the encryption key string formed the encryption key to be plugged into the algorithm).

Summary of the Invention

In general terms, the present invention calls for the recovery of encrypted sensitive data to require the involvement not only of a first trusted authority competent in respect of the accreditation of professionals, but also of an organisation engaging the professional and a second trusted authority competent in respect of the accreditation of organisations.

More particularly, according to a first aspect of the present invention, there is provided a method of recovering target data provided in encrypted form to a party as part of a data set with which first and second trusted authorities are associated in a non-subvertible manner, the method comprising:

- providing a first element to the party after the first trusted authority has verified that a specific individual is a professional accredited with it;
 - providing a second element to the party after both the second trusted authority has verified that a particular organisation is accredited with it, and said particular organisation has verified that said specific individual is engaged by it; and
 - the party using both said elements to recover the target data in clear;
- at least one of the particular organisation and the first trusted authority ensuring that its verification is for said party as said specific individual before providing the corresponding element.

In one embodiment both the particular organisation and the first trusted authority use the authenticated identity of the party for the specific individual in respect of which they carry out their respective verifications. In another embodiment, the data set identifies said specific individual and one or both of the particular organisation and the first trusted

authority check that the authenticated identity of the party corresponds to the specific individual identified in the data set.

Advantageously, the method involves the use of Identifier-Based Encryption (IBE). In one preferred embodiment, the data set comprises a first item encrypted using a first IBE encryption key that identifies said specific individual, and public data of the first trusted authority; and a second item encrypted using a second IBE encryption key that identifies a specific organisation, and public data of the second trusted authority. In this case, the second trusted authority verifies that the said particular organisation is the specific organisation identified in the second encryption key as well it as being an organisation accredited with the second trusted authority.

The use of the public data of the first and second trusted authorities in encrypting the first and second items provides a non-subvertible link between the data set and the trust authorities as these authorities must be contacted for the corresponding decryption keys. However, it may be noted that the data provider may opt to use the same first and second encryption keys when encrypting the first and second items of different data sets in which case provision can be made for caching of the corresponding decryption keys, thereby obviating the need for the trusted authorities to be contacted each time target data is provided to the party.

According to a second aspect of the present invention, there is provided a secure data-provision method comprising providing target data from a data provider to a party purporting to be a specific, professionally-accredited, individual engaged by a specific accredited organisation, the target data being provided in encrypted form as part of a data set that comprises:

- a first item encrypted, according to an Identifier-Based Encryption, IBE, scheme, using both a first encryption key that identifies said specific individual, and public data of a first trusted authority competent in respect of professional accreditations;
- and

- a second item encrypted according to an IBE scheme, using both a second encryption key that identifies said specific organisation, and public data of a second trusted authority competent in respect of accreditations of organisations;

recovery of the target data in clear requiring decryption of both the first and second items.

5

According to a third aspect of the present invention, there is provided a system for recovering target data provided in encrypted form to a party as part of a data set with which first and second trusted authorities are associated in a non-subvertible manner, the system comprising:

- 10 - a first computing entity, associated with the first trusted authority, for providing a first element to the party after verifying that a specific individual is a professional accredited with it;
- a second computing entity associated with the second trusted authority;
- a third computing entity, associated with a particular organisation, for providing a
- 15 second element to the party after the second computing entity has verified that said particular organisation is accredited with it, and the third computing entity has verified that said specific individual is engaged by it; and
- a fourth computing entity, associated with said party, for decrypting the target data using the first and second elements;
- 20 at least one of the first and third computing entities being arranged to ensure that its verification is for said party as said specific individual before providing the corresponding element to the party.

According to a fourth aspect of the present invention, there is provided apparatus for the

25 secure provision of target data to a party purporting to be a specific, professionally-accredited, individual engaged by a specific accredited organisation, the apparatus comprising an encryption subsystem for generating a data set including the target data in encrypted form, the encryption subsystem comprising:

- first encryption means for encrypting a first item, according to an Identifier-Based
- 30 Encryption, IBE, scheme, using both a first encryption key that identifies said specific individual, and public data of a first trusted authority competent in respect of professional accreditations;

- second encryption means for encrypting a second item, according to an IBE scheme, using both a second encryption key that identifies said specific organisation, and public data of a second trusted authority competent in respect of accreditations of organisations; and
- 5 - means for forming the data set using at least the encrypted first and second items; the recovery of the target data in clear requiring decryption of both the first and second items.

The present invention also envisages user computing devices for use by professionals in
10 recovering encrypted target data.

Brief Description of the Drawings

Embodiments of the invention will now be described, by way of non-limiting example,
15 with reference to the accompanying diagrammatic drawings, in which:

- . Figure 1 is a diagram illustrating the operation of a prior art encryption schema known as Identifier-Based Cryptography;
- . Figure 2 is a diagram showing the general arrangement of entities involved in the embodiments described with respect to Figures 3 to 8;
- 20 . Figure 3 is a diagram of a first specific embodiment of the present invention;
- . Figure 4 is a diagram of a second specific embodiment of the present invention;
- . Figure 5 is a diagram of a third specific embodiment of the present invention;
- . Figure 6 is a diagram of a fourth specific embodiment of the present invention;
- . Figure 7 is a diagram of a first variant of the Figure 5 embodiment; and
- 25 . Figure 8 is a diagram of a second variant of the Figure 5 embodiment.

Best Mode of Carrying Out the Invention

The embodiments of the invention to be described hereinafter are all placed in a medical context with a requesting party only being able to obtain access to a patient record if the
30 party is a medical professional (for example, a doctor or paramedic) accredited with a medical professional trusted authority and engaged by a medical organisation (such as a hospital) accredited with a medical organisation trusted authority. However, it is to be

understood that these embodiments can also be applied in other fields beyond the medical world.

Figure 2 illustrates the general arrangement of the entities involved in all the embodiments described below. More particularly, this arrangement comprises a first computing entity 20 (such as a personal digital assistant) associated with a requesting party wishing to receive a patient record; a second computing entity 30 associated with a patient record storage service; a third computing entity 40 associated with a medical professional trusted authority competent in respect of the accreditation of medical professionals (that is, trusted as authoritative concerning the accreditation of such professionals); a fourth computing entity 45 associated with a medical organisation trusted authority competent in respect of the accreditation of medical organisations; and a fifth computing entity 50 associated with a particular medical organisation. The computing entities 20, 30, 40, 45 and 50 are typically based around general-purpose processors executing stored programs. The computing entities 20, 30, 40, 45 and 50 inter-communicate as needed (see arrows 55-58) via, for example, the internet or other network though it is also possible that at least some of the entities actually reside on the same computing platform. As will be described below, at least certain of the inter-entity communications are arranged to take place securely with the communicating parties authenticating each other; to this end, the entities are equipped with suitable communication modules well understood by persons skilled in the art.

In the following, references to the requesting party, patient record storage service, the medical professional trusted authority, the medical organisation trusted authority, and the medical organisation are generally used interchangeably with references to their respective computing entities 20, 30, 40, 45 and 50. Furthermore, for convenience the terms "patient record storage service" and "trusted authority" are abbreviated to "PRSS" and "TA" respectively.

In functional terms, the requesting-party entity 20 comprises a communications module 23 for communicating with the entities 30, 40 and 50, a control module 21 for controlling the general operation of the entity 20 and for providing a user interface and at least short-term

storage, and a cryptographic module 22 for executing certain cryptographic functions that vary between the embodiments to be described below.

5 The PRSS entity 30 comprises a communications module 34 for communicating with the requesting party entity 20 (and possibly also with the entities 40 and 45), a control module 31 for controlling the general operation of the entity 30, a database 32 for holding patient records, and a cryptographic module 33 for executing certain cryptographic functions that also vary between the embodiments to be described below.

10 The medical professional TA entity 40 comprises a communications module 44 for communicating with the requesting party entity 20 (and possibly also with the entity 30), a control module 41 for controlling the general operation of the entity 40, a database 42 for holding medical professional accreditation data, and a cryptographic module 43 for executing certain cryptographic functions.

15

The medical organisation TA entity 45 comprises a communications module 49 for communicating with the medical organisation entity 50 (and possibly also with the entity 30), a control module 46 for controlling the general operation of the entity 45, a database 47 for holding medical organisation accreditation data, and a cryptographic module 48 for
20 executing certain cryptographic functions.

The medical organisation entity 50 comprises a communications module 54 for communicating with the requesting party entity 20 and the medical organisation TA 45, a control module 51 for controlling the general operation of the entity 50, a database 52 for
25 holding data about medical professionals engaged by the organisation including their data access authorisation levels (in particular, whether they are authorised to access patient records), and a cryptographic module 53 for executing certain cryptographic functions.

The specific embodiments now to be described all employ Identifier-Based Cryptography
30 (in the present case, the QR IBC method) to enable the PRSS entity 30 to specify conditions to be met by parties wishing to access patient records provided by the entity 30.

More particularly, the TAs 40 and 45 have respective IBE public data $N1$ and $N2$ and corresponding respective IBE private data $p1, q1$ and $p2, q2$ used in forming their public data. The PRSS entity 30 knows the public data $N1$ and $N2$ of the two TAs (for example, as a result of the latter each publishing its public data in a certificate digitally signed using
 5 a locally-held private key of a public/private key pair associated with the trusted authority).

When the requesting party 20 wants to access a patient record, it makes a request (arrow 55) to the PRSS entity 30 in which it not only identifies the patient concerned, but also identifies both itself (by name or, preferably, by another identifier such as a public key of
 10 an asymmetric public/private key pair the private key of which is held by the party 20), and the medical organisation for which the party 20 is currently working (again, either by name or by another identifier such as the public key of an asymmetric public/private key pair the private key of which is held by the organisation).

15 The PRSS entity 30 responds to the request by the party 20 by encrypting the requested patient record (referred to herein as the "target record" or, more generally, the "target data") and providing it (arrow 55) to the party 20 as part of a data set that comprises encrypted first and second items. The first data-set item is IBE encrypted using the party's supplied identity as an IBE encryption key and the public data $N1$ of the medical
 20 professional TA 40. The second data-set item is IBE encrypted using the supplied organisation identity as an IBE encryption key and the public data $N2$ of the medical organisation TA 45. To recover the target patient record in clear, it is necessary to decrypt both the first data-set item and the second data-set item and this requires a first IBE decryption key provided by the medical professional TA 40 and a second decryption key
 25 provided by the medical organisation TA 45.

As will become apparent hereinafter, the composition of the data set of which the encrypted target patient record forms a part varies from embodiment to embodiment as does the relationship between the first and second data-set items and the target patient
 30 record (indeed, in one embodiment, the first data-set item is the target patient record).

- The party entity 20 on receiving the data set including the encrypted target record, seeks to obtain the first decryption key from the medical professional TA 40 and in doing so provides the related encryption key to the TA 40. The TA 40 only returns the decryption key if it is satisfied that the individual identified in the encryption key is a professional accredited with it as indicated by the data it holds in its database 42; the TA 40 may also require the party to prove that they are this identified individual. In certain embodiments, the TA 40 may be arranged to receive, decrypt and return the first data-set item rather than providing the first decryption key to the party 20.
- 10 The party entity 20 also requests (arrow 57) the second decryption key from the medical organisation entity 50, providing the latter with the encryption key that identifies the organisation indicated to the PRSS entity by the party 20. The organisation 50, either before or after carrying out certain checks to be described, asks (arrow 58) the medical organisation TA 45 to provide the second decryption key on the basis of the supplied
- 15 encryption key. The TA 45 only supplies the requested key if it is satisfied that the requesting organisation is the same organisation as identified in the encryption key and that the organisation is accredited with it according to the data held in the database 47. Assuming that the TA 45 provides the second decryption key to the organisation 50; and provided this entity 50 is satisfied that the party 20 (or, in certain embodiments, the
- 20 individual identified by the party to the PRSS entity 30), is engaged by the organisation with appropriate data access authority as indicated by data in the database 52, the organisation returns (arrow 57) the second decryption key, or the second data-set item decrypted using this key, to the party 20.
- 25 The final recovery of the target patient record takes place at the party entity 20. This recovery is only possible if the party 20 is a medical professional accredited with the medical professional TA 40 and is engaged by a medical organisation accredited with the medical organisation TA 45. However, it may be noted that the PRSS entity 30 may use the same encryption keys when encrypting the first and second items of data sets associated
- 30 with different record requests by the party 20; in this case, the corresponding decryption keys may be cached by the entities that carry out IBE decryption, thereby obviating the

need for the TAs 40, 45 to be contacted each time a target record is provided to the party 20.

So far as the IBE cryptographic processes are concerned, the correspondence between the 5 entities of Figure 2 and those of Figure 1 will be apparent to persons skilled in the art.

Specific IBE-based embodiments will now be described with reference to Figures 3 to 8. In these Figures, the PRSS entity 30 and the initial request from the party entity to the PRSS entity are not depicted and consideration of these embodiments starts with the data set 10 returned by the PRSS entity 30, it being appreciated that this data set has been generated by the functional elements 31 and 33 of the entity 30 using the data supplied to it in the initial request, the TA public data N1 and N2, and the patient record extracted from the database 32.

15 Considering first the Figure 3 embodiment, the data set returned by the PRSS entity 30 comprises:

- a first IBE encryption key K1 comprising an identification of an individual (the party 20) purporting to be a medical professional (MP);
- a second IBE encryption key K2 comprising an identification of an organization 20 purported by the party 20 as being an accredited medical organization ("MO") by which the party is engaged;
- an encrypted first item formed by the IBE encryption of the target patient record "PR" using the first encryption key K1 and the public data N1 of the medical professional TA 40 – this is represented by the expression $E\langle K1, N1; PR \rangle$ where $E\langle \rangle$ denotes that 25 the element appearing in the brackets after the semi colon has been encrypted using the element or elements appearing before the semi colon;
- an encrypted second item formed by the IBE encryption of the first item using the second encryption key K2 and the public data N2 of the medical organization TA 45 – this is represented by the expression $E\langle K2, N2; E\langle K1, N1; PR \rangle \rangle$.

30 It will be appreciated that the encrypted first item (that is, the second item) does not appear explicitly in the data set but only in its further encrypted form as the encrypted second item.

The entity 20 now establishes a secure authenticated communication channel 100 with the medical professional TA 40 and requests (arrow 60) the IBE decryption key K3 corresponding to IBE encryption key K1, this latter being passed in the request to the entity 40. The entity 40 first checks (process 61) that the requesting party, as established by authentication when setting up channel 100, is the same as the medical professional (MP) identified in the encryption key K1. If this check is passed, the entity 40 then checks (process 62) that the party/MP is accredited as a medical professional with the entity 40. The checks 61 and 62 can, in fact, be carried out simultaneously or in reverse order. Only if both these checks are passed does the entity 40 proceed with the generation (process 63) of the decryption key K3 by using the encryption key K1 and the private data $p1, q1$ of the entity 40. The decryption key K3 is then returned (arrow 64) over the channel 100 to the party entity 20. The key K3 could have been generated in advance of, or in parallel with, the checks 61 and 62 being carried out – what is important is that the key K3 is not returned until the checks have been passed.

The entity 20 also establishes a secure authenticated communication channel 101 with the medical organization entity 50 and passes it both the IBE encryption key K2 (arrow 65) and the encrypted second item (arrow 66). The entity 50 first checks (process 67) that the party 20, authenticated during set up of the channel 101, is an individual engaged by it with authority to access patient records. “Engaged” can either be taken as currently engaged over a sustained period (for example, of weeks, months, years or for an unspecified duration), or be taken to be actually on duty for the organization at the current instance. The entity 50 may (or may not) also check that it is the organization identified in the encryption key K2. If the or each of these checks is passed, the medical organization entity 50 sets up a secure authenticated communication channel 102 with the medical organization TA entity 45 and passes it (arrow 68) the encryption key K2 with a request for the corresponding decryption key K4.

The TA entity 45 first checks (process 69) that the requesting organization, as authenticated during set up of channel 102, is the organization identified in the encryption key K2. If this check is passed, the entity 45 then checks (process 70) that the organisation

is accredited as a medical organisation with the entity 45. The checks 69 and 70 can be carried out simultaneously or in reverse order. Only if both these checks are passed does the entity 45 proceed with the generation (process 71) of the decryption key K4 by using the encryption key K2 and the private data $p2, q2$ of the entity 45. The decryption key K4 is then returned (arrow 72) over the channel 102 to the medical organisation entity 50. The key K4 could have been generated in advance of, or in parallel with, the checks 69 and 70 being carried out – what is important is that the key K4 is not returned until the checks have been passed.

- 10 On receiving the decryption key K4, the medical organisation entity 50 uses it to decrypt (process 73) the encrypted second item. The second item $E\langle K1, N1; PR \rangle$ is then passed back (arrow 74) over the secure channel 101 to the party entity 20.

The party entity 20 finally recovers the target patient record in clear by using the decryption
15 key K3 to decrypt the second item (process 80).

In the embodiments of Figures 4, 5 and 6 to be described below, the roles of the two TA entities 40 and 45 and of the medical organisation entity 50 are substantially the same as
20 for the Figure 3 embodiment, namely:

- the medical professional TA entity 40, after carrying out checks 61 and 62, provides the party entity 20 with an IBE decryption key K2 for decrypting a first item encrypted with a first IBE encryption key K1,
- the medical organisation TA entity 45, after carrying out checks 69 and 70, provides
25 the medical organisation entity 20 with an IBE decryption key K4 for decrypting a second item encrypted with a second IBE encryption key K2, and
- the medical organisation entity 50, after carrying out check 67, decrypts the second item using the key K4 and returns the decrypted second item to the party entity 20.

The differences between the embodiments of Figures 3 to 6 lie in the contents of the data
30 sets provided by the PRSS entity 30 and how these contents are used to recover the target patient record in clear.

Considering next the embodiment of Figure 4, in this embodiment the data set provided by the PRSS entity 30 comprises:

- identification of an individual (the party 20) purporting to be a medical professional (MP), this identification being that used by the PRSS entity 30, in combination with a
5 nonce (random number), for a first IBE encryption key K1;
- a second IBE encryption key K2 comprising an identification of an organization purported by the party 20 as being an accredited medical organization ("MO") by which the party is engaged;
- an encrypted first item $E\langle K1, N1; PR \rangle$ formed by the IBE encryption of the target
10 patient data record "PR" using the first encryption key K1 and the public data N1 of the medical professional TA 40;
- an encrypted second item $E\langle K2, N2; Nonce \rangle$ formed by the IBE encryption of the nonce used in the first encryption key K1, using the second encryption key K2 and the public data N2 of the medical organization TA 45.

15 In this embodiment, the party entity 20 first obtains the decrypted second item (the nonce used in the encryption key K1) from the medical organisation entity 50 and then uses this nonce, together with the medical professional identifier provided by the PRSS entity 30, to re-form (process 81) the encryption key K1. The re-formed key K1 is passed to the entity 40 to obtain the corresponding decryption key K3 which is then used to decrypt (process
20 82) the encrypted first item and recover the target patient record in clear.

Considering the embodiment of Figure 5, in this embodiment the data set provided by the PRSS entity 30 comprises:

- a first IBE encryption key K1 comprising identification of an individual (the party 20)
25 purporting to be a medical professional (MP);
- a second IBE encryption key K2 comprising an identification of an organization purported by the party 20 as being an accredited medical organization ("MO") by which the party is engaged;
- the target patient record encrypted using a symmetric key S;
- an encrypted first item $E\langle K1, N1; A \rangle$ formed by the IBE encryption of a first part A of
30 the symmetric key S using the first encryption key K1 and the public data N1 of the medical professional TA 40;

- an encrypted second item $E\langle K2, N2; B \rangle$ formed by the IBE encryption of a second part B of the symmetric key S using the second encryption key K2 and the public data N2 of the medical organization TA 45.

In this embodiment, the party entity 20 obtains the decryption key K3 from the medical professional TA entity 40 and uses it to decrypt (process 83) the encrypted first item to provide the first part A of the symmetric key S. The entity 20 also obtains the decrypted second item (the second part B of the symmetric key S) from the medical organisation entity 50. The party entity 20 then combines A and B (process 84) to re-form the symmetric key S which it thereafter uses to decrypt the target patient record (process 85).

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Rather than the symmetric key S being simply split into two parts A and B and re-formed by concatenation of A and B, a more complex relationship between S, A and B is preferred that avoids disclosure of A or B providing any information about S. By way of example, A and B could be created first and then S derived as a hash of A and B, i.e. $S = \text{Hash}(A, B)$. An alternative approach would be to use Shamir's security sharing.

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Considering the embodiment of Figure 6, in this embodiment the data set provided by the PRSS entity 30 comprises:

- a first IBE encryption key K1 comprising identification of an individual (the party 20) purporting to be a medical professional (MP);
- a second IBE encryption key K2 comprising an identification of an organization purported by the party 20 as being an accredited medical organization ("MO") by which the party is engaged;
- the target patient record encrypted using a first symmetric key S1;
- an encrypted first item $E\langle K1, N1; E\langle S2; S1 \rangle \rangle$ formed by the IBE encryption of a component using the first encryption key K1 and the public data N1 of the medical professional TA 40, the component concerned being the first symmetric key S1 encrypted using a second symmetric key S2;
- an encrypted second item $E\langle K2, N2; S2 \rangle$ formed by the IBE encryption of the second symmetric key S2 using the second encryption key K2 and the public data N2 of the medical organization TA 45.

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In this embodiment, the party entity 20 obtains the decryption key K3 from the medical professional TA entity 40 and uses it to decrypt (process 86) the encrypted first item to provide the component $E\langle S2; S1 \rangle$. The entity 20 also obtains the decrypted second item (the second symmetric key S2) from the medical organisation entity 50 which it then uses
 5 to decrypt (process 87) the component $E\langle S2; S1 \rangle$ and obtain the first symmetric key S1. The party entity 20 then uses the first symmetric key S1 to decrypt the encrypted target patient record (process 88).

It will be appreciated that the data set provided by the PRSS entity 20 may take many other
 10 forms without affecting the roles played by the entities 40, 45 and 50. Other variants that modify the operation of the entities 40, 45 and 50 are also possible. For example, in the embodiments of Figures 3 to 6 if the PRSS entity 30 is arranged to use the same first encryption key K1 when encrypting the first items of successive data sets provided to the same party (that is, data sets provided in response to successive patient record requests
 15 from the party 20), it is possible to arrange for the corresponding decryption key K3 to be cached by the entity 20 when first obtained from the TA entity 40 for use in decrypting the first item of the first of these data sets, and then have the entity 20 re-use this key K3 from cache for the other data sets that have their first items encrypted with same encryption key as the first item for which the decryption key was obtained. In a similar manner, if the
 20 PRSS entity 30 uses the same encryption key K2 for encrypting the second items of successive data sets, the medical organisation entity 50 can cache the corresponding decryption key K4 when first obtained from the TA entity 45 and subsequently use the key from cache. Of course, re-use of either decryption key from cache means that the checks carried out by the corresponding TA entity are avoided; it is therefore preferable for the
 25 PRSS entity 30 to limit either the number of times or the period over which it re-uses the same encryption key; for example, the PRSS entity 30 may be arranged to change the first encryption key K1 once a month and to change the second encryption key K3 daily. Changing an encryption key whilst retaining the identification of the medical professional or medical organisation identified in the key is readily done by including a fresh nonce
 30 each time the key is to be changed; every time a new nonce is included in an encryption, the corresponding TA entity 40, 45 must be involved to generate the corresponding decryption key thereby resulting in the checks associated with the TA entity to be effected.

Of course, the encryption keys K1 and K2 could be changed at every usage; however, this may not be desirable where speed of retrieval of a patient record is important as in emergency medical situations.

5 Another variant generally applicable to all the embodiments of Figures 3 to 6 is for the TA entity 45, rather than passing the decryption key K3 to the party entity 20, to use this key itself to decrypt the encrypted first item and then provide the first item back to the party entity 20. Similarly, the TA entity 45, rather than passing the decryption key K4 to the organisation entity 50, can use this key itself to decrypt the encrypted second item and then
 10 provide the second item back to the organisation entity 20. It is also possible to arrange for the medical organisation entity 50 to pass back to the party entity 20 the element it receives from the TA entity 45, whether this be the decryption key K4 or the recovered second element, without using the element itself. These variants are not, however, preferred and in any event the final decryption of the patient record should normally be restricted to being
 15 done at the party entity 20.

Figure 7 shows a further variant here applied to the Figure 5 embodiment but also applicable to the embodiments of Figures 3, 4 and 6. In the Figure 7 variant the second encryption key K2 comprises not only identification of a medical organisation, but also the
 20 identification of a medical professional contained in the first encryption key K1. This enables the medical organisation entity 50 to check (process 90) that the medical professional checked out by the medical professional TA entity 45 is also engaged by the organisation 50 with authority to access patient records. Whilst the medical organisation entity 50 preferably also checks that the party with which it is communicating is the
 25 medical professional identified in the second encryption key, this is not essential because the medical organisation can be sure that only the medical professional identified in the encryption keys will be able to receive the first decryption key K3 from the medical professional TA entity 40. The role of the medical organisation 50 is thus simply to provide the key K4 if the medical professional identified in the first encryption key to the medical
 30 professional TA entity 40 is engaged by an organisation accredited with the medical organisation entity 45. Indeed, provision of the key K4 to the party 20 can be done over an insecure, unauthenticated, channel 104 as the key K4 will only be of value to a party that

also possesses the first decryption key K3; provision over a secure channel is, however, preferred.

Figure 8 illustrates another variant applied to the Figure 5 embodiment but also applicable to the embodiments of Figures 3, 4 and 6. The Figure 8 variant, like that of Figure 7, is based on the second encryption key including not only identification of a medical organisation, but also identification of the same medical professional as identified in the first encryption key. In this variant, the medical organisation entity 50 both checks that the party 20 is the medical professional identified in the second encryption key K2 (process 91) and checks that this professional is engaged by the organisation 50 with authority to access patient records (process 90). Since the decrypted second item will only be available to the party who is the medical professional identified in the encryption keys, the medical professional TA entity 40 need only check that the professional identified in the first encryption key K1 is accredited with it and return the first decryption key K3 to the requesting party; the TA entity 40 need neither check that the party 20 is the identified medical professional nor provide the key K3 in a secure manner (in Figure 8 the entities 20 and 40 are depicted as communicating via an insecure, unauthenticated, channel 105). Of course, in reality communication between the entities 20 and 40 is preferably by a secure channel and the entity 40 preferably does check that the party 20 is the professional identified in the first encryption key.

It will also be appreciated that many other variants are possible to the above described embodiments of the invention. For example, the data set provided by the PRSS entity 30 need not all be provided to the party 20 but components of it could be passed to the entities 40 and 50 directly for their use. Similarly, since the party entity 20 may well be connected to a network run by the medical organisation 50, the latter can be arranged to intercept the data set and copy or strip out the components it needs before passing on the data set, or the remainder of it, to the party 20.

In the above-described embodiments, no restrictions have been placed on the professional or organisation identified by the party to the PRSS entity 30 when requesting a patient record. However, the PRSS entity 30 can be arranged to authenticate the party 30 and use

the authenticated identity of the party in the first encryption key K1. Furthermore, the party entity 20 can take the form of a trusted computing platform provided by the organisation 50 and adapted to reliably provide the PRSS entity 30 with the identifications of both the party 20 and the organisation 50. One suitable form of trusted platform is specified in "TCPA -

5 Trusted Computing Platform Alliance Main Specification v1.1" www.trustedcomputing.org, 2001 and described in the book "trusted computing platforms - tcpa technology in context"; Pearson (editor); Prentice Hall; ISBN 0-13-009220-7". The computing entity 50 is, preferably, also a trusted computing platform the status of which is verifiable by the TA entity 45.

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Where the identification of the professional and organisation to the PRSS entity 30 is not controlled or checked, then there may be little to be gained by using these identifications in the IBE encryption keys in which case the checks carried out by the entities 40, 45 and 50 are simply that the party 20 itself is a medical professional accredited with the medical

15 professional TA entity 40 and engaged by an organisation 50 which is accredited with the medical organisation TA entity 45.

The present invention is not limited to the QR IBE method used in the above-described

20 embodiments and other, analogous, cryptographic methods can be used such as IBE methods based on Weil or Tate pairings.

Furthermore, embodiments of the invention based on cryptographic techniques other than IBE are also possible. For example, in variant of the Figure 5 embodiment, the PRSS entity

25 30 encrypts the quantity A and an identifier of a medical professional using a public key of a first public/private asymmetric key pair the private key of which is held by the medical professional TA entity 40. Similarly, the PRSS entity 30 also encrypts the quantity B and an identifier of a medical organisation using a public key of a second public/private asymmetric key pair the private key of which is held by the medical organisation TA entity

30 45. In this case, the TA entities 40 and 45 first use their respective private keys to decrypt the data encrypted using their public keys, and then carry out their checks 61,62; 69,70 as described above for the Figure 5 embodiment before passing the quantities A and B

respectively to the party 20 and the medical organisation 50. The organisation 50 passes on the quantity B to the party 20 after carrying out its check 67.

5 The variants discussed above in relation to the IBE embodiments (including, in particular, those of Figures 7 and 8) are generally also applicable to the non-IBE embodiments; for example, the encrypted data provided to the TA entities in the example described in the foregoing paragraph need not include identifications of a medical professional or medical organisation.

CLAIMS

1. A method of recovering target data provided in encrypted form to a party as part of a data set with which first and second trusted authorities are associated in a non-subvertible manner, the method comprising:

- providing a first element to the party after the first trusted authority has verified that a specific individual is a professional accredited with it;
- providing a second element to the party after both the second trusted authority has verified that a particular organisation is accredited with it, and said particular organisation has verified that said specific individual is engaged by it; and
- the party using both said elements to recover the target data in clear;

at least one of the particular organisation and the first trusted authority ensuring that its verification is for said party as said specific individual before providing the corresponding element.

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2. A method according to claim 1, wherein said data set comprises an encrypted first item decryptable using a first decryption key obtainable from the first trusted authority and an encrypted second item decryptable using a second decryption key obtainable from the second trusted authority, both items requiring decryption for the target data to be recovered in clear; the method comprising:

- (a) - providing the party, in a secure manner, with the first element from the first trusted authority only if the latter is satisfied that the party is a professional accredited with this authority, the first element being one of the first decryption key and the first item recovered with this key;
- (b) - providing said particular organisation, in a secure manner, with the second decryption key, or the second item decrypted with this key, from the second trusted authority only if the latter is satisfied that the particular organisation is accredited with this authority;
- (c) - providing the party with the second element from said particular organisation only if the organisation is satisfied that the party is engaged by the organisation with authority to access the target data, the second element being one of the second decryption key and the second item recovered with this key; and

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- (d) - recovering the target data at the party by using the decrypted first and second items, decryption of the target data being effected at the party.

3. A method according to claim 2, wherein the particular organisation provides its output
5 to said party in a secure manner.

4. A method according to claim 1, wherein said data set comprises an encrypted first item decryptable using a first decryption key obtainable from the first trusted authority and an encrypted second item decryptable using a second decryption key obtainable from the
10 second trusted authority, both items requiring decryption for the target data to be recovered in clear and the data set identifying said specific individual; the method comprising:

(a) - providing the party with the first element from the first trusted authority only if the latter is satisfied that said specific individual, as identified in the data set, is a professional accredited with this authority, the first element being one of the first
15 decryption key and the first item recovered with this key;

(b) - providing said particular organisation, in a secure manner, with the second decryption key, or the second item decrypted with this key, from the second trusted authority only if the latter is satisfied that the organisation is accredited with this authority;

20 (c) - providing the party with the second element, from said particular organisation only if the latter is satisfied that said specific individual is engaged by the organisation with authority to access the target data, the second element being one of the second decryption key and the second item recovered with this key; and

(d) - recovering the target data at the party by using the decrypted first and second items,
25 decryption of the target data being effected at the party;

at least one of the first trusted authority and said particular organisation providing its associated element to said party in a secure manner and only after being satisfied that the party is the said specific individual identified in the data set.

30 5. A method according to any one of claims 2 to 4, wherein:

- the first item has been encrypted, according to an Identifier-Based Encryption, IBE, scheme, using a first encryption key that identifies said specific individual, and

public data of the first trusted authority; the first decryption key being a key generated by the first trusted authority using the first encryption key and private data used in forming said public data; and

- the second item has been encrypted, according to an IBE scheme, using a second encryption key that identifies a specific organisation, and public data of the second trusted authority; the second decryption key being a key generated by the second trusted authority using the second encryption key and private data used in forming said public data; the second decryption key, or the second item recovered with this key, only be provided to said particular organisation upon the second trusted authority being satisfied that said particular organisation is the specific organisation identified in the second encryption key as well it as being an organisation accredited with the second trusted authority.

6. A method according to claim 5, wherein the first item comprises the target data, and the second item comprises the encrypted first item; the second item being recovered using the second decryption key and then being subject to decryption using the first decryption key to recover the target data.

7. A method according to claim 5, wherein the first item comprises the target data, the second item comprises a nonce, and the first encryption key comprises, in combination, an identifier of said specific individual and said nonce; the second item being recovered using the second decryption key to provide said nonce which is then combined with the identifier of said specific individual to form the first encryption key, this key thereafter being provided to the first trusted authority in order to enable the latter to provide the first decryption key for use at said party to decrypt the first item and thereby recover the target data.

8. A method according to claim 5, wherein the first item comprises first data and the second item comprises second data, the data set further comprising said target data encrypted using a symmetric key that can be formed by using both said first and second data; the first and second items being independently recovered to provide the first and

second data which are then used at said party to form said symmetric key, this key then being used to decrypt the target data.

9. A method according to claim 5, wherein the data set comprises, in addition to said first
5 and second items, said target data encrypted using a first symmetric key, the second item
comprising a second symmetric key, and the first item comprising the first symmetric key
encrypted using the second symmetric key; the first and second items being independently
recovered to provide the encrypted first symmetric key and the second symmetric key, and
10 the second symmetric key then being used at said party to decrypt the encrypted first
symmetric key after which the first symmetric key is used to decrypt the encrypted target
data.

10. A method according to claim 5, wherein the party is provided with multiple different
encrypted target datas in respective data sets, the first item of each data set being encrypted
15 with a first encryption key that comprises, in addition to identification of said specific
individual, a random element; the party, in recovering each target data, obtaining from the
first trusted authority, the corresponding first decryption key or the corresponding first item
recovered using this key.

20 11. A method according to claim 5, wherein the party is provided with multiple different
encrypted target datas in respective data sets, the first item of each data set being encrypted
with the same first encryption key; the party, after having obtained the corresponding first
decryption key from the first trusted authority in the course of recovering the target data of
one data set, caching the first decryption key and re-using it from cache for recovering the
25 target datas of subsequent data sets that have first items encrypted using the same first
encryption key as that for which the cached first decryption key was obtained.

12. A method according to claim 5, wherein the party is provided with multiple different
encrypted target datas in respective data sets, the second item of each data set being
30 encrypted with a second encryption key that comprises, in addition to identification of said
specific organisation, a random element; the requesting organisation obtaining from the

second trusted authority, for each different second encryption key used, the corresponding second decryption key or the corresponding second item recovered using this key.

13. A method according to claim 5, wherein the party is provided with multiple different
5 encrypted target datas in respective data sets, the second item of each data set being encrypted with the same second encryption key; the requesting organisation after having obtained the corresponding second decryption key from the second trusted authority in respect of one data set, caching the second decryption key and re-using it from cache for subsequent data sets that have second items encrypted using the same second encryption
10 key as that for which the cached second decryption key was obtained, the requesting organisation decrypting the second item itself and passing the recovered second item to said party.

14. A secure data-provision method comprising providing target data from a data provider
15 to a party purporting to be a specific, professionally-accredited, individual engaged by a specific accredited organisation, the target data being provided in encrypted form as part of a data set that comprises:

- a first item encrypted, according to an Identifier-Based Encryption, IBE, scheme, using both a first encryption key that identifies said specific individual, and public
20 data of a first trusted authority competent in respect of professional accreditations; and
 - a second item encrypted according to an IBE scheme, using both a second encryption key that identifies said specific organisation, and public data of a second trusted authority competent in respect of accreditations of organisations;
- 25 recovery of the target data in clear requiring decryption of both the first and second items.

15. A method according to claim 14, wherein the first item comprises the target data, and the second item comprises the encrypted first item.

30 16. A method according to claim 14, wherein the first item comprises the target data, and the second item comprises a nonce; the first encryption key comprising, in combination, an identifier of said specific individual and said nonce.

17. A method according to claim 14, wherein the first item comprises first data, and the second item comprises second data; the data set further comprising said target data encrypted using a symmetric key that can be formed by using both said first and second
5 data.

18. A method according to claim 14, wherein the data set comprises, in addition to said first and second items, said target data encrypted using a first symmetric key, the second item comprising a second symmetric key, and the first item comprising the first symmetric
10 key encrypted using the second symmetric key.

19. A secure data-provision method comprising providing target data from a data provider to a party purporting to be a specific, professionally-accredited, individual engaged by a specific accredited organisation, the target data being provided in encrypted form as part of
15 a data set that comprises:

- a first item encrypted using both a first encryption key that identifies said specific individual, and public data of a first trusted authority competent in respect of professional accreditations; and
- a second item encrypted using both a second encryption key that identifies said
20 specific organisation, and public data of a second trusted authority competent in respect of accreditations of organisations;

recovery of the target data in clear requiring decryption of both the first and second items.

20. A system for recovering target data provided in encrypted form to a party as part of a
25 data set with which first and second trusted authorities are associated in a non-subvertible manner, the system comprising:

- a first computing entity, associated with the first trusted authority, for providing a first element to the party after verifying that a specific individual is a professional accredited with it;
- 30 - a second computing entity associated with the second trusted authority;
- a third computing entity, associated with a particular organisation, for providing a second element to the party after the second computing entity has verified that said

particular organisation is accredited with it, and the third computing entity has verified that said specific individual is engaged by it; and

- a fourth computing entity, associated with said party, for decrypting the target data using the first and second elements;
- 5 at least one of the first and third computing entities being arranged to ensure that its verification is for said party as said specific individual before providing the corresponding element to the party.

21. A system according to claim 20, wherein:

- 10 - said data set comprises an encrypted first item decryptable using a first decryption key obtainable from the first trusted authority and an encrypted second item decryptable using a second decryption key obtainable from the second trusted authority, both items requiring decryption for the target data to be recovered in clear;
- the first computing entity is arranged to provide said party, in a secure manner, with
15 the first element only if it is satisfied that the party is a professional accredited with the first trusted authority, the first element being one of the first decryption key and the first item recovered with this key;
- the second computing entity is arranged to provide said particular organisation, in a secure manner, with the second decryption key, or the second item decrypted with
20 this key, only if it is satisfied that the particular organisation is accredited with the second trusted authority;
- the third computing entity is arranged to provide said party with the second element only if it is satisfied that the party is engaged by said particular organisation with authority to access the target data, the second element being one of the second
25 decryption key and the second item recovered with this key; and
- the fourth computing entity is arranged to recover the target data by using the decrypted first and second items.

22. A system according to claim 21, wherein the third computing entity is arranged to
30 provide its output to said party in a secure manner.

23. A system according to claim 20, wherein:

- said data set comprises an encrypted first item decryptable using a first decryption key obtainable from the first trusted authority and an encrypted second item decryptable using a second decryption key obtainable from the second trusted authority, both items requiring decryption for the target data to be recovered in clear and the data set identifying said specific individual;
 - the first computing entity is arranged to provide said party with the first element only if it is satisfied that said specific individual, as identified in the data set, is a professional accredited with the first trusted authority, the first element being one of the first decryption key and the first item recovered with this key;
 - the second computing entity is arranged to provide said particular organisation, in a secure manner, with the second decryption key, or the second item decrypted with this key, only if it is satisfied that the organisation is accredited with the second trusted authority;
 - the third computing entity is arranged to provide the party with the second element only if it is satisfied that said specific individual is engaged by the organisation with authority to access the target data, the second element being one of the second decryption key and the second item recovered with this key; and
 - the fourth computing entity is arranged to recovering the target data by using the decrypted first and second items;
- at least one of the first and third computing entities being arranged to provide the corresponding element to said party in a secure manner and only after being satisfied that the party is the said specific individual identified in the data set.

24. A system according to any one of claims 21 to 23, wherein:

- the first item has been encrypted, according to an Identifier-Based Encryption, IBE, scheme, using a first encryption key that identifies said specific individual, and public data of the first trusted authority; the first computing entity being arranged to generate the first decryption key using the first encryption key and private data used in forming said public data; and
- the second item has been encrypted, according to an IBE scheme, using a second encryption key that identifies a specific organisation, and public data of the second trusted authority; the second computing entity being arranged to generate the second

decryption key using the second encryption key and private data used in forming said public data, and the second computing entity being further arranged to provide the second decryption key, or the second item recovered using this key, to said particular organisation upon being satisfied that said particular organisation is the specific organisation identified in the second encryption key as well it as being an organisation accredited with the second trusted authority.

25. A system according to claim 24, wherein the first item comprises the target data, and the second item comprises the encrypted first item; the fourth computing entity being arranged to:

- recover the second item, if not provided to it in decrypted form by the third computing entity, by using the second decryption key obtained from the third computing entity, and
- subject the second item to decryption, using the first decryption key obtained from the first computing entity, to recover the target data.

26. A system according to claim 24, wherein the first item comprises the target data, the second item comprises a nonce, and the first encryption key comprises, in combination, an identifier of said specific individual and said nonce; the fourth computing entity being arranged to:

- recover the second item, if not provided to it in decrypted form by the third computing entity, by using the second decryption key obtained from the third computing entity,
- combine the nonce that formed the second item with the identifier of said specific individual in order to form the first encryption key,
- provide the first encryption key to the first computing entity to obtain the first decryption key, and
- use the first decryption key obtained from the first computing entity to decrypt the first item and thereby recover the target data.

27. A system according to claim 24, wherein the first item comprises first data and the second item comprises second data, the data set further comprising said target data

encrypted using a symmetric key that can be formed by using both said first and second data; the fourth computing entity being arranged to

- recover the first data, if not provided to it in decrypted form by the first computing entity, by using the first decryption key obtained from the first computing entity,
- 5 - recover the second data, if not provided to it in decrypted form by the third computing entity, by using the second decryption key obtained from the third computing entity,
- use the first data and the second data to form said symmetric key, and
- use the symmetric key to decrypt the target data.

10

28. A system according to claim 24, wherein the data set comprises, in addition to said first and second items, said target data encrypted using a first symmetric key, the second item comprising a second symmetric key, and the first item comprising the first symmetric key encrypted using the second symmetric key; the fourth computing entity being arranged

15 to:

- recover the first item, if not provided to it in decrypted form by the first computing entity, by using the first decryption key obtained from the first computing entity,
- recover the second item, if not provided to it in decrypted form by the third computing entity, by using the second decryption key obtained from the third computing entity,
- 20 - use the second symmetric key that formed the second item to decrypt the encrypted first symmetric key that formed the first item, and
- use the first symmetric key to decrypt the encrypted target data.

20

25 29. Apparatus for the secure provision of target data to a party purporting to be a specific, professionally-accredited, individual engaged by a specific accredited organisation, the apparatus comprising an encryption subsystem for generating a data set including the target data in encrypted form, the encryption subsystem comprising:

- first encryption means for encrypting a first item, according to an Identifier-Based Encryption, IBE, scheme, using both a first encryption key that identifies said specific individual, and public data of a first trusted authority competent in respect of professional accreditations;

30

- second encryption means for encrypting a second item, according to an IBE scheme, using both a second encryption key that identifies said specific organisation, and public data of a second trusted authority competent in respect of accreditations of organisations; and
- 5 - means for forming the data set using at least the encrypted first and second items; the recovery of the target data in clear requiring decryption of both the first and second items.
30. Apparatus according to claim 29, wherein the first item comprises the target data, and
- 10 the second item comprises the encrypted first item.
31. Apparatus according to claim 29, wherein the first item comprises the target data, and the second item comprises a nonce; the first encryption key comprising, in combination, an identifier of said specific individual and said nonce.
- 15
32. Apparatus according to claim 29, wherein the first item comprises first data, and the second item comprises second data; the data set further comprising said target data encrypted using a symmetric key that can be formed by using both said first and second data.
- 20
33. Apparatus according to claim 29, wherein the data set comprises, in addition to said first and second items, said target data encrypted using a first symmetric key, the second item comprising a second symmetric key, and the first item comprising the first symmetric key encrypted using the second symmetric key.
- 25
34. A computing entity for recovering target data provided in encrypted form as part of an data set that comprises first and second encrypted items both of which must be decrypted to recover the target data, the first item being encrypted using both a first encryption key that identifies a specific individual and first public data, and the second item being encrypted
- 30 using both a second encryption key that identifies a specific organisation, and second public data; the entity comprising:

- first means for requesting either a first decryption key corresponding to the first encryption key, or the first item in decrypted form, from a first trusted authority which is competent in respect of the accreditation of professionals and holds first private data used in deriving the first public data, the first means being arranged to provide the first encryption key to the first trusted authority when making its request and being further arranged to authenticate the entity with the first trusted authority and to receive the first decryption key, or the first item, securely from the first trusted authority;
- second means for requesting either a second decryption key corresponding to the second encryption key, or the second item in decrypted form, from an organisation accredited by a second trusted authority which holds second private data used in deriving the second public data, the second means being arranged to provide the second encryption key to the organisation when making its request and being further arranged to authenticate the entity with the organisation and receive the second decryption key, or the second item, from the organisation;
- third means for using the first decryption key, or the first item, provided by the first trusted authority and the second decryption key, or the second item, provided by the organisation, to recover the target data.

35. A computing entity according to claim 30, wherein the second means is arranged to receive the second decryption key, or the second item, securely from the organisation.

36. A computing entity for recovering target data provided in encrypted form as part of an data set that comprises first and second encrypted items both of which must be decrypted to recover the target data, the first item being encrypted using both a first encryption key that identifies a specific individual and first public data, and the second item being encrypted using both a second encryption key that identifies a specific organisation and said specific individual, and second public data; the entity comprising:

- first means for requesting either a first decryption key corresponding to the first encryption key, or the first item in decrypted form, from a first trusted authority which is competent in respect of the accreditation of professionals and holds first private data used in deriving the first public data, the first means being arranged to

provide the first encryption key, or the first item, to the first trusted authority when making its request;

- second means for requesting either a second decryption key corresponding to the second encryption key, or the second item in decrypted form, from an organisation accredited by a second trusted authority which holds second private data used in deriving the second public data, the second means being arranged to provide the second encryption key to the organisation when making its request; and

- third means for using the first decryption key, or the first item, provided by the first trusted authority and the second decryption key, or the second item, provided by the organisation, to recover the target data;

at least one of the first means and the second means being arranged to authenticate the entity to the first trusted authority or said organisation as the case may be and to receive input therefrom in a secure manner.

37. A computing entity according to any one of claims 34 to 36, wherein the first item comprises the target data, and the second item comprises the encrypted first item; the third means being arranged to:

- recover the second item, if not provided to the second means in decrypted form by the organisation, by using the second decryption key obtained from the organisation, and
- subject the second item to decryption, using the first decryption key obtained from the first trusted authority, to recover the target data.

38. A computing entity according to any one of claims 34 to 36, wherein the first item comprises the target data, the second item comprises a nonce, and the first encryption key comprises, in combination, an identifier of said specific individual and said nonce; the third means being arranged to:

- recover the second item, if not provided to the second means in decrypted form by the organisation, by using the second decryption key obtained from the organisation,
- combine the nonce that formed the second item with the identifier of said specific individual in order to form the first encryption key to be provided by the first means to the first trusted authority, and

- use the first decryption key obtained from the first trusted authority to decrypt the first item and thereby recover the target data.

39. A computing entity according to any one of claims 34 to 36, wherein the first item
5 comprises first data and the second item comprises second data, the data set further comprising said target data encrypted using a symmetric key that can be formed by using both said first and second data; the third means being arranged to

- recover the first data, if not provided to the first means by the first trusted authority, by using the first decryption key obtained from the first trusted authority,
- 10 - recover the second data, if not provided to the second means in decrypted form by the organisation, by using the second decryption key obtained from the organisation,
- use the first data and the second data to form said symmetric key, and
- use the symmetric key to decrypt the target data.

15 40. A computing entity according to any one of claims 34 to 36, wherein the data set comprises, in addition to said first and second items, said target data encrypted using a first symmetric key, the second item comprising a second symmetric key, and the first item comprising the first symmetric key encrypted using the second symmetric key; the third means being arranged to:

- 20 - recover the first item, if not provided to the first means by the first trusted authority, by using the first decryption key obtained from the first trusted authority,
- recover the second item, if not provided to the second means in decrypted form by the organisation, by using the second decryption key obtained from the organisation,
- use the second symmetric key that formed the second item to decrypt the encrypted
- 25 first symmetric key that formed the first item, and
- use the first symmetric key to decrypt the encrypted target data.

ABSTRACT

5 **Secure Data Provision Method and Apparatus and**
 Data Recovery Method and System.

10 To control access to target data whilst relieving the data provider (30) of policing obligations, the data provider (30) provides the target data in encrypted form to a requesting party (20) as part of a data set with which first and second trusted authorities (40,45) are associated in a non-subvertible manner. Recovery of the target data in clear by the party (20) requires the first trusted authority (40) to verify that a specific individual is a professional accredited with it, the second trusted authority (45) to verify that a particular organisation (50) is accredited with it, the particular organisation (50) to verify that the specific individual is engaged by it, and at least one of the particular organisation (50) and the first trusted authority (40) to verify that the party (20) is the specific individual. Various ways of encrypting the target data are provided, the preferred ways being based on Identifier-Based Encryption schemas.

20

(Figure 2)

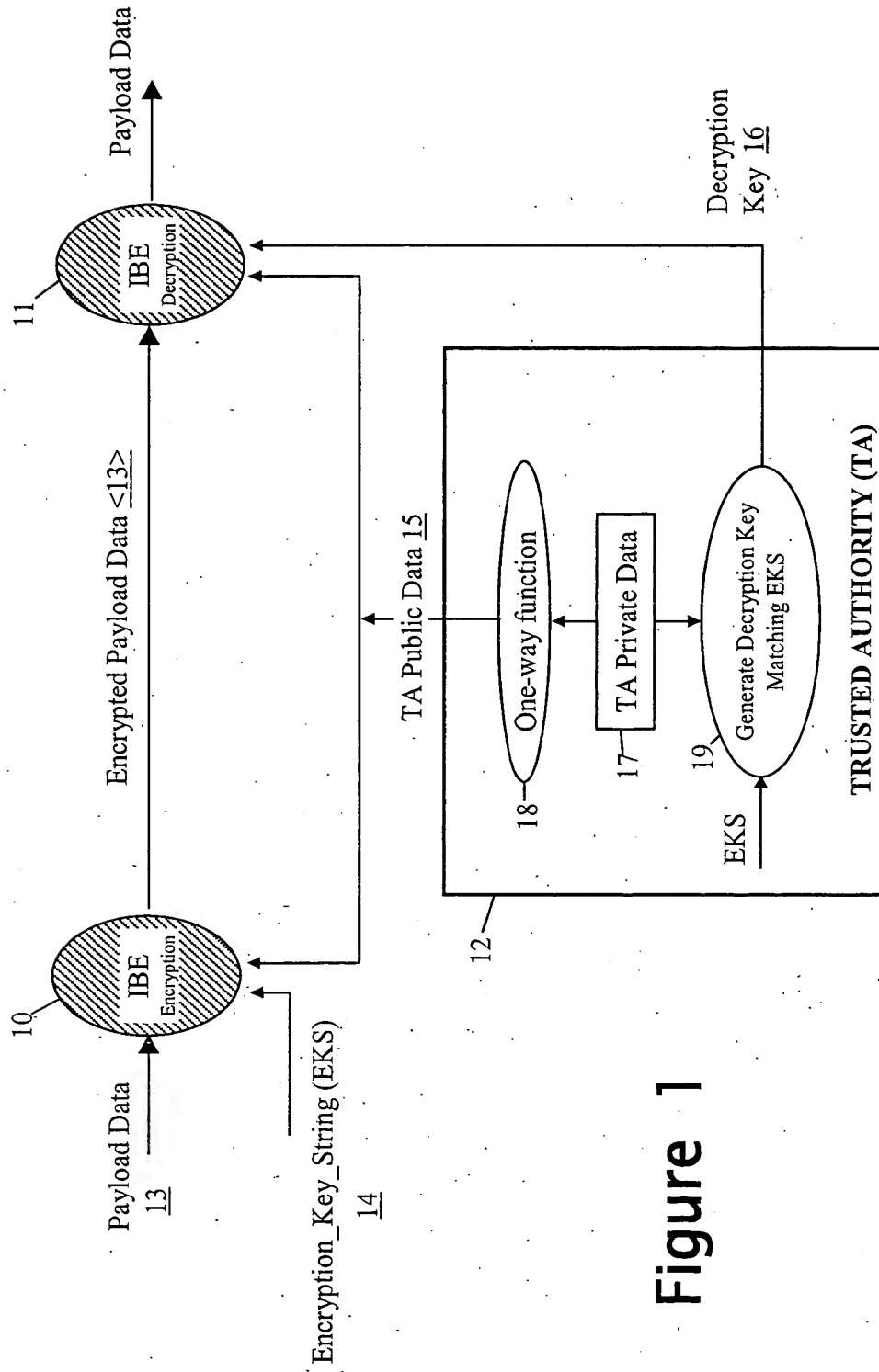


Figure 1

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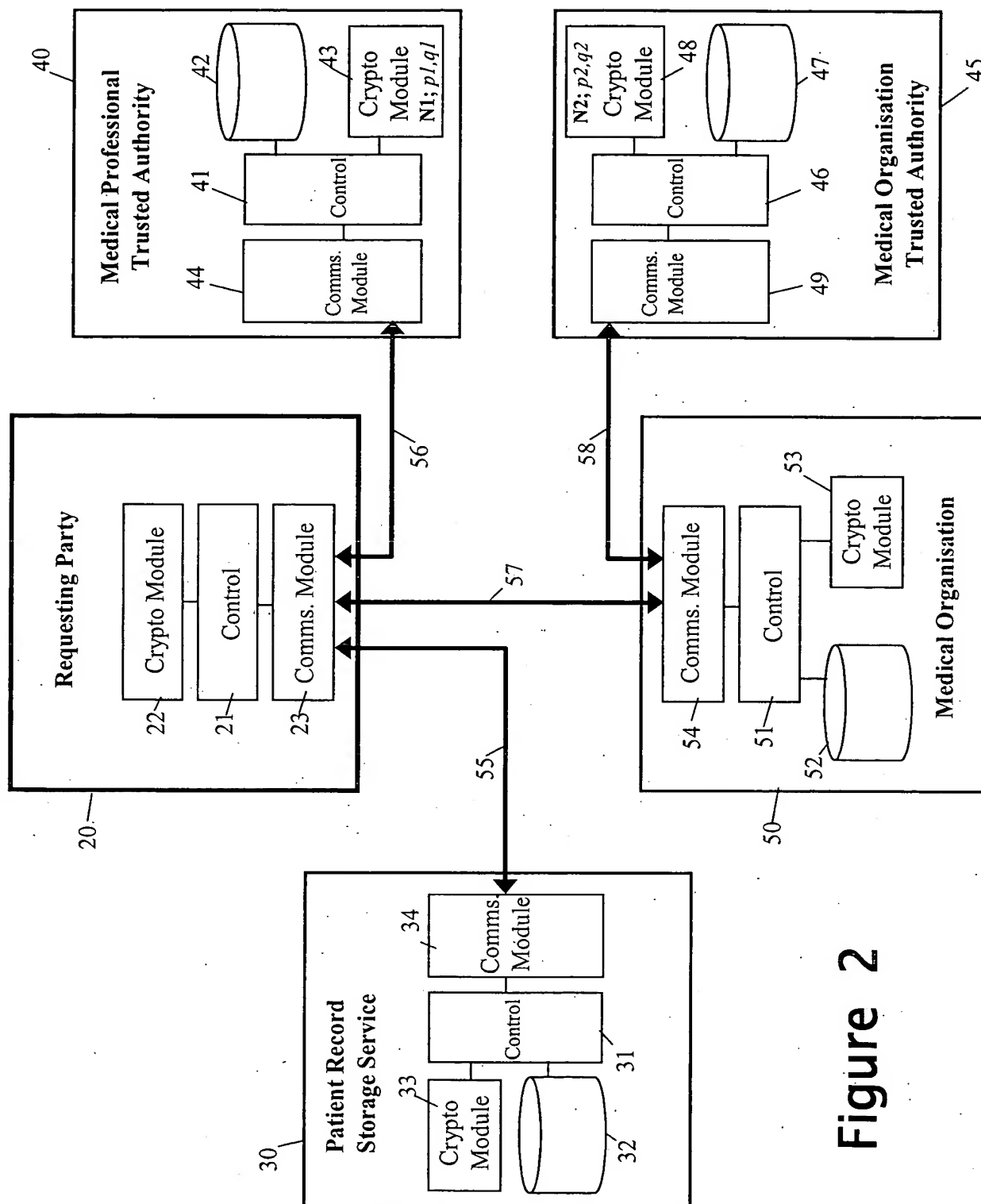


Figure 2

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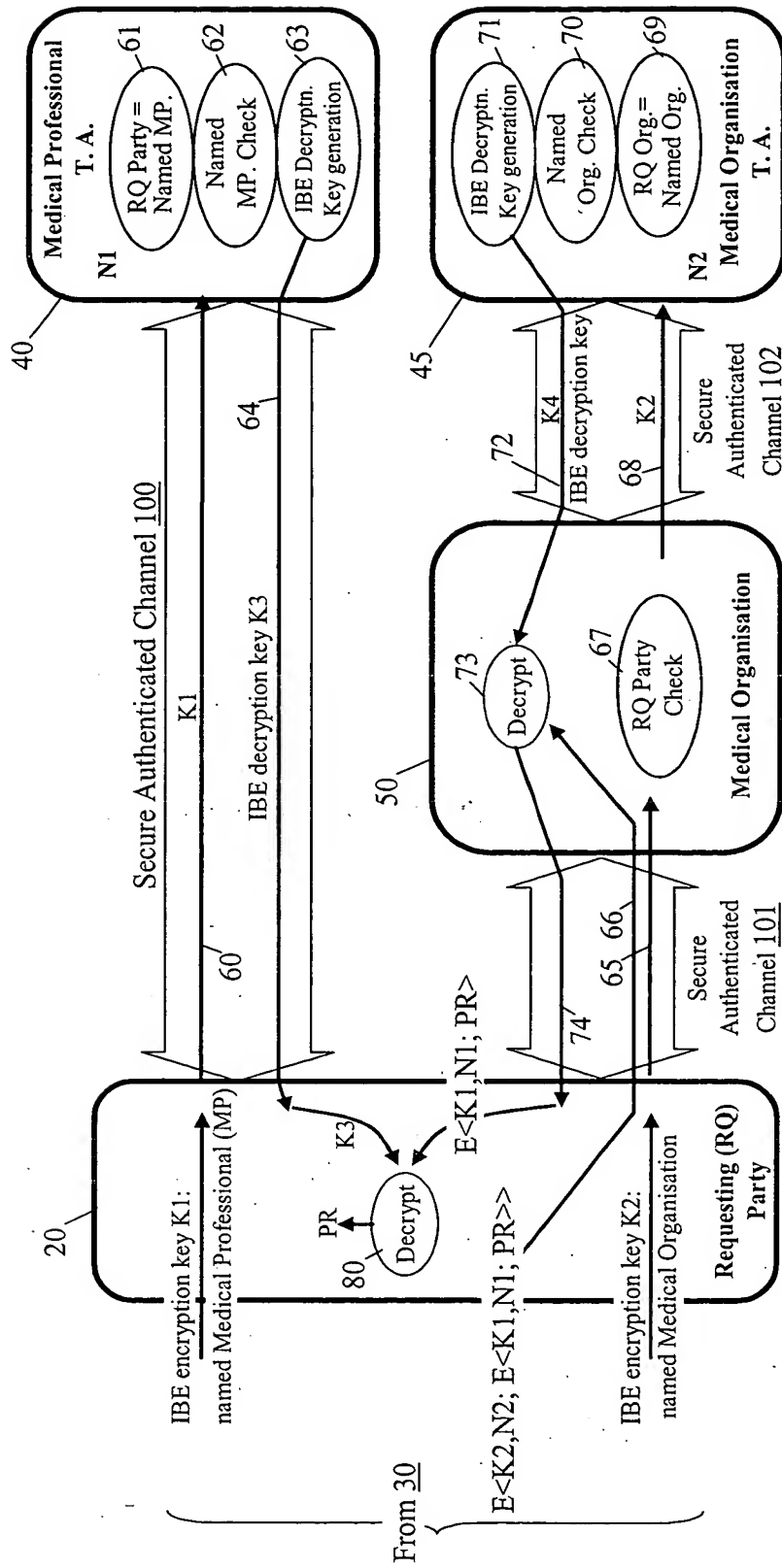


Figure 3

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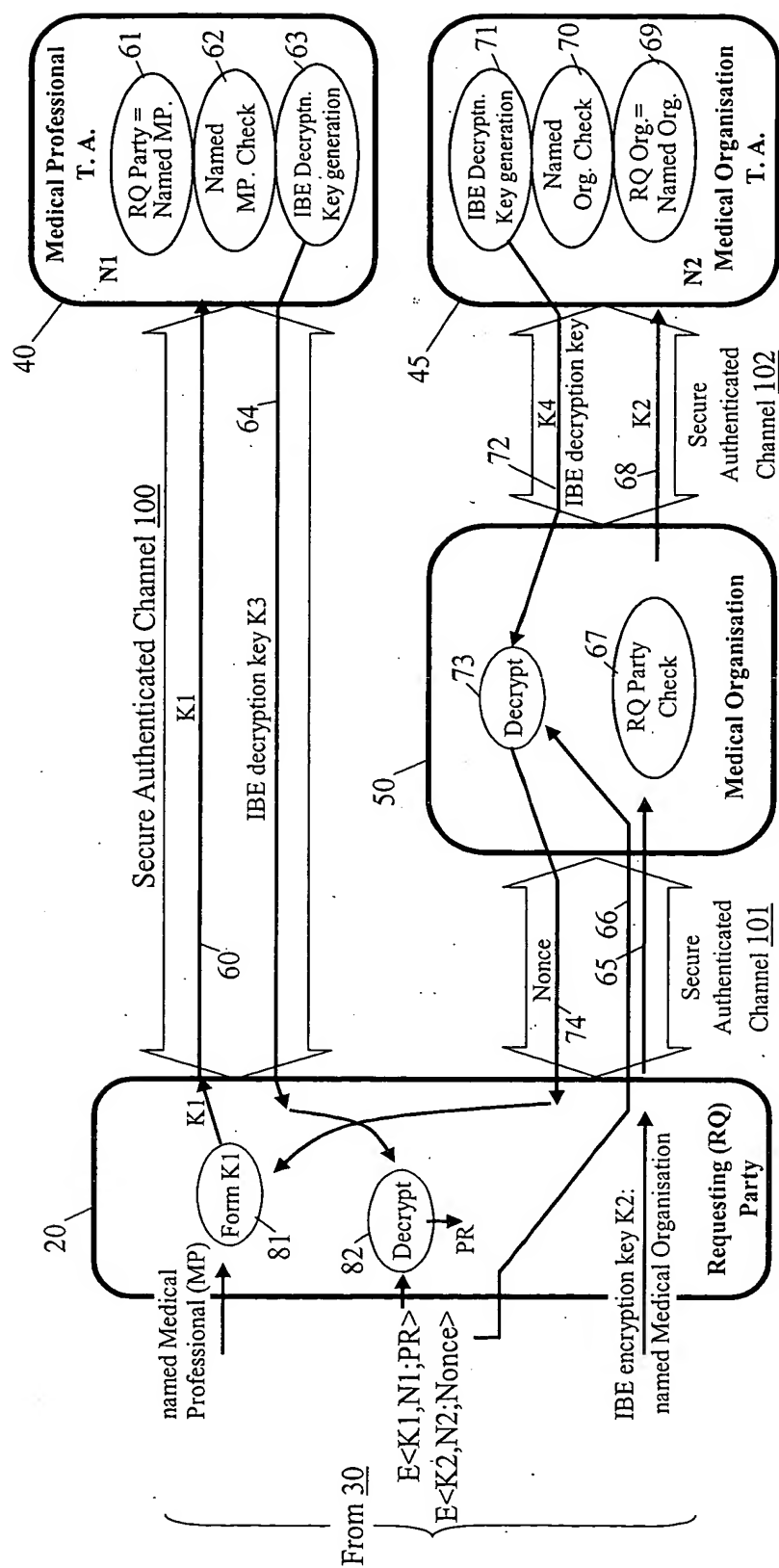


Figure 4

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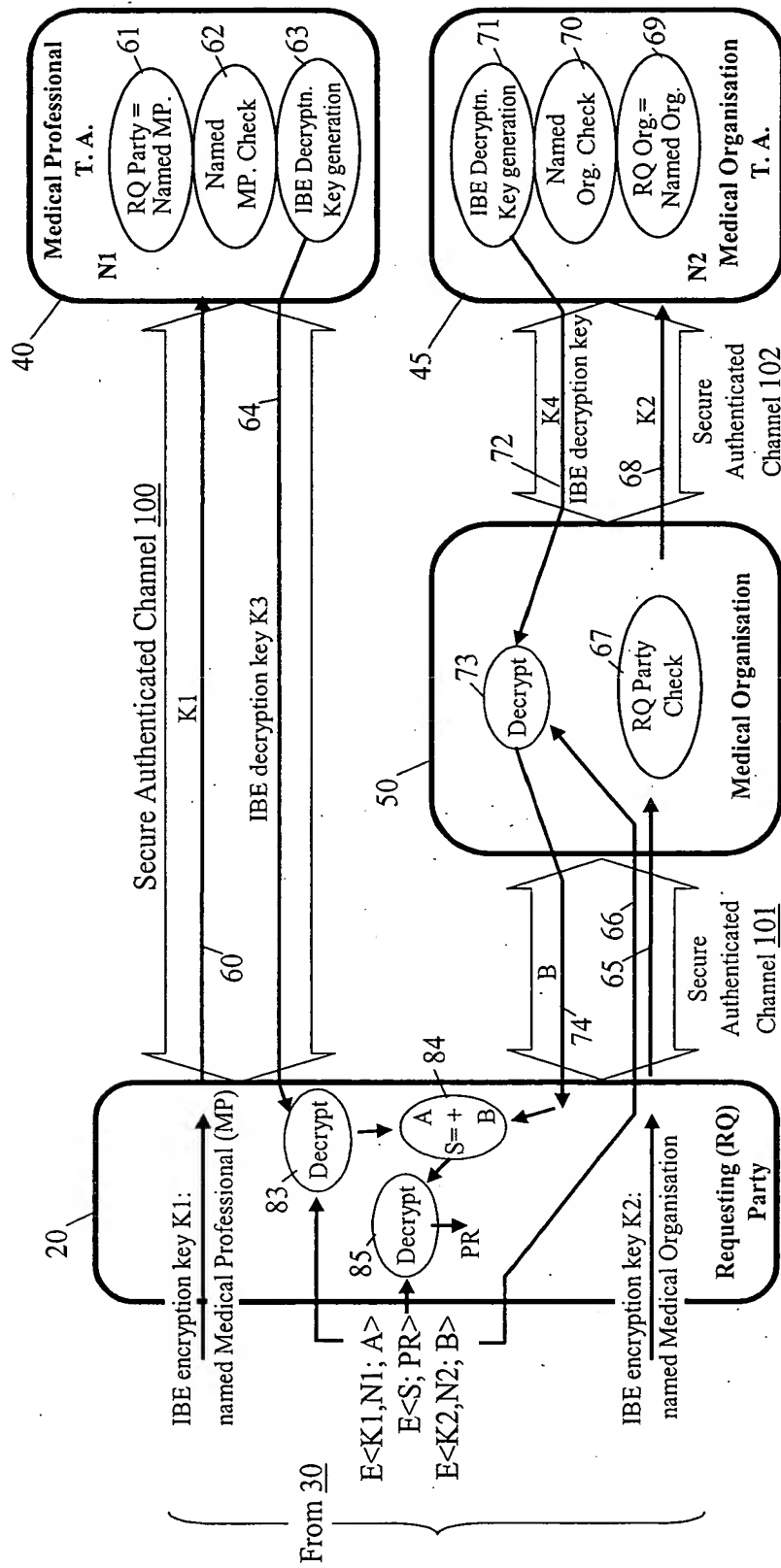


Figure 5

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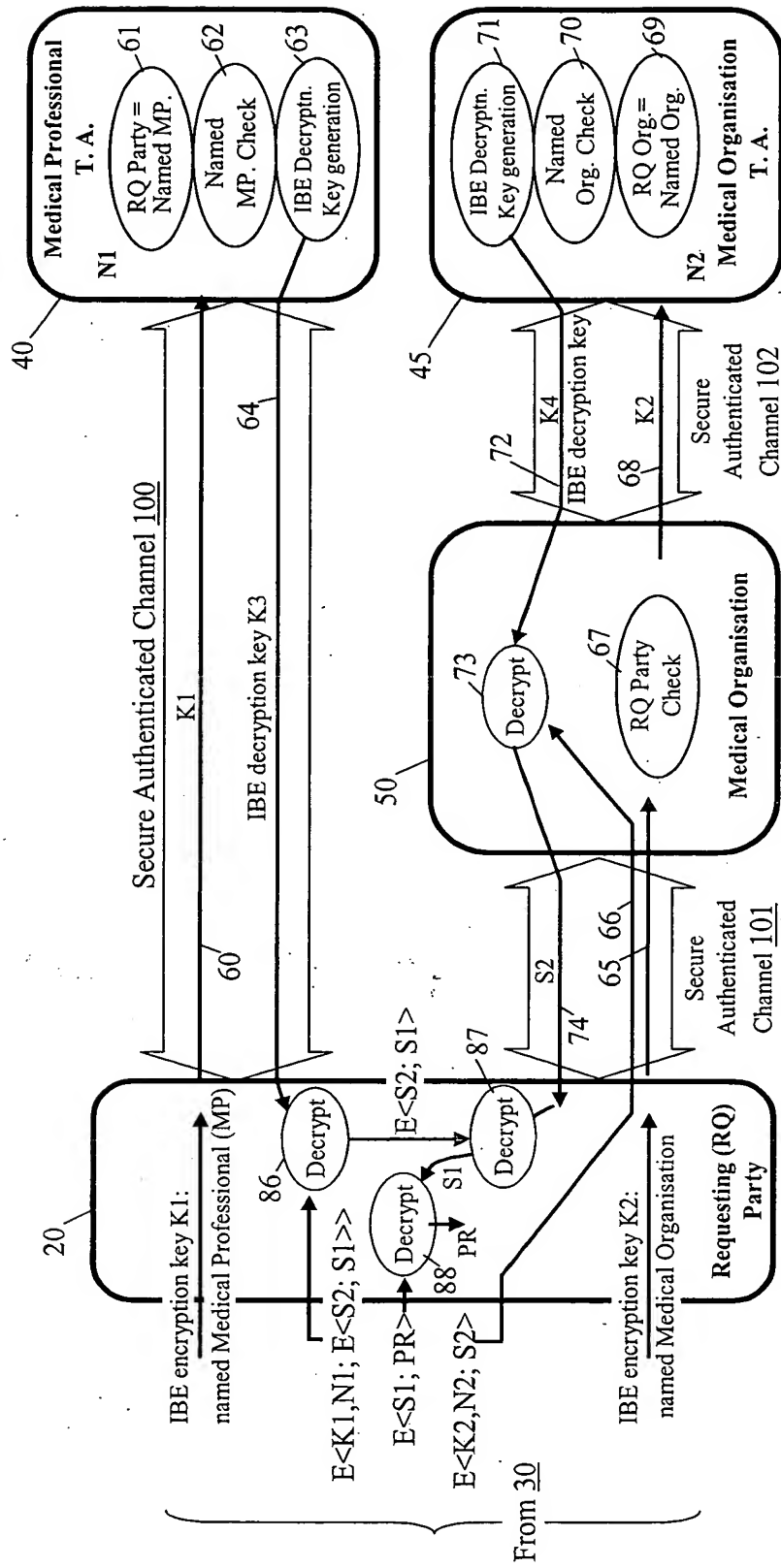


Figure 6

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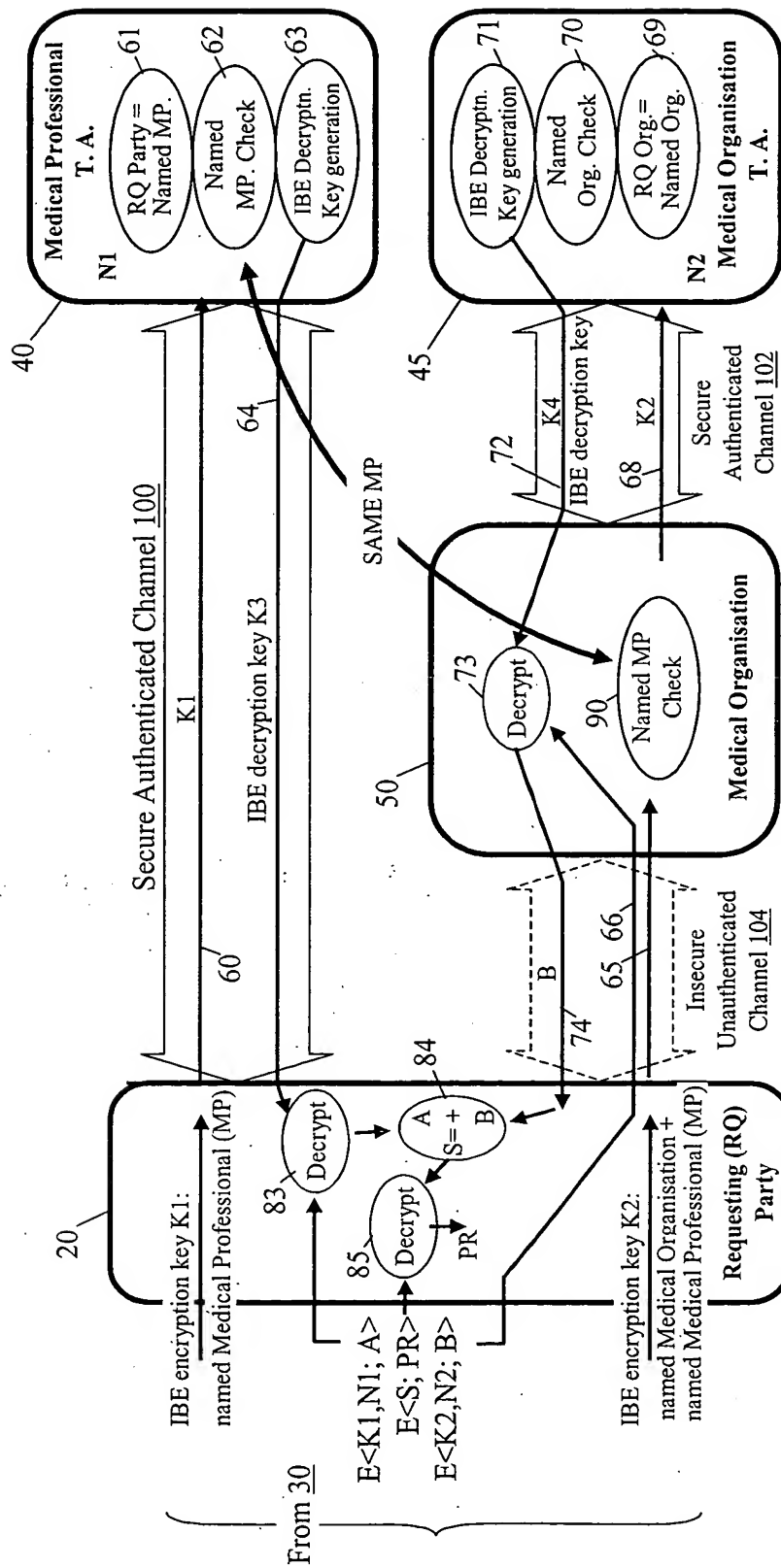


Figure 7

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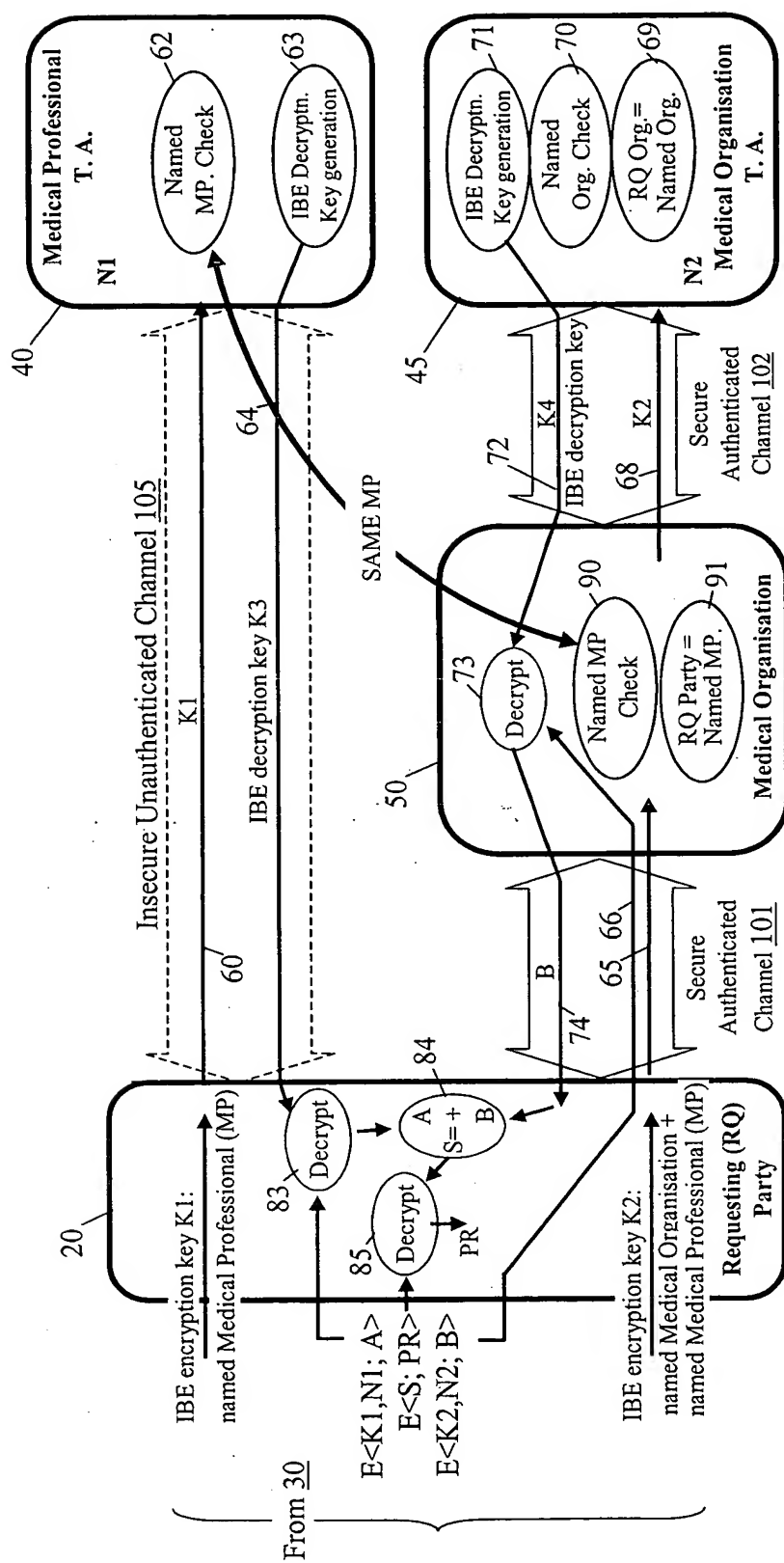


Figure 8

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